

Geometry of log-concave Ensembles of random matrices

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Basic definitions and Notation

Let $X = (X(1), \dots, X(N))$ be a random vector in \mathbb{R}^N with full dimensional support. We say that the distribution of X is

- *log-concave*, if X has density of the form $e^{-h(x)}$ with $h: \mathbb{R}^N \rightarrow (-\infty, \infty]$ convex;
- *isotropic*, if $\mathbb{E}X_i = 0$ and $\mathbb{E}X_i X_j = \delta_{i,j}$.

For $x, y \in \mathbb{R}^N$ we put

- $|x| = \|x\|_2 = \left(\sum_{i=1}^N x_i^2 \right)^{1/2}$
- $\langle x, y \rangle$ is the inner product
- For $\emptyset \neq I \subset \{1, \dots, N\}$, by P_I we denote the coordinate projection onto $\{y \in \mathbb{R}^N: \text{supp}(y) \subset I\}$.

Examples: log-concave random vectors

1. Let $K \subset \mathbb{R}^n$ be a convex body (= compact convex, with non-empty interior) (symmetric means $-K = K$).

X a random vector uniformly distributed in K . Then the corresponding probability measure on \mathbb{R}^n

$$\mu_K(A) = \frac{|K \cap A|}{|K|}$$

is log-concave (by **Brunn-Minkowski**).

Moreover, for every convex body K there exists an affine map T such that μ_{TK} is isotropic.

2. The Gaussian vector $G = (g_1, \dots, g_n)$, where g_i 's have $\mathcal{N}(0, 1)$ distribution, is isotropic and log-concave.

3. Similarly the vector $X = (\xi_1, \dots, \xi_n)$, where ξ_i 's are independent with the exponential distribution (i.e., with density $f(t) = \frac{1}{\sqrt{2}} \exp(-\sqrt{2}|t|)$, for $t \in \mathbb{R}$) is isotropic and log-concave.

Let n, N be integers; $X \in \mathbb{R}^N$ a random vector, X_1, \dots, X_n independent, distributed as X .

A is an $n \times N$ matrix which has X_i 's as rows

$$A = \begin{pmatrix} X_1 \\ X_2 \\ \dots \\ X_n \end{pmatrix} = \begin{pmatrix} \dots\dots\dots \\ \dots\dots\dots \\ \dots\dots\dots \\ \dots\dots\dots \end{pmatrix}$$

For $J \subset \{1, \dots, n\}$, $I \subset \{1, \dots, N\}$, let $A(J, I)$ be the sub-matrix with rows from J and columns from I . For $k \leq n$, $m \leq N$, let

$$A_{k,m} = \sup_{J,I} \|A(J, I)\|,$$

over all $|J| \leq k$, $|I| \leq m$. This is maximum of norms of sub-matrices of k columns and m rows (the operator norms from ℓ_2^N to ℓ_2^n).

Question: Upper bounds for $A_{k,m}$ with high probability

Example: Norms of sub-matrices, uniform version of P.

Earlier case: $m = N$, i.e., $A_{k,N}$. We select k rows and take *all* columns. Solved by ALPT (JAMS 2010); connected with length of good approximations of the covariance matrices

Case $k = 1$: For any J with $|J| = 1$, and I with $|I| \leq m$, sub-matrix of one row has a form

$$A(J, I) = \left(P_I X \right), \quad \text{so} \quad A_{1,m} = \sup_{|I| \leq m} |P_I X|.$$

A bound for $A_{1,m}$ will imply a uniform bound for $P_I X$, for $t \geq 1$.

$$\sup_{|I|=m} |P_I X| \leq \text{bound for } A_{1,m}, \quad (*)$$

with high probability. Compare with Paouris' theorem for a *fixed* I with $|I| \leq m$: for $s \geq 1$,

$$\mathbb{P} \left(|P_I X| \leq C s \sqrt{N} \right) \geq 1 - \exp \left(- s \sqrt{N} \right).$$

However the complexity of the family of subsets is $\binom{N}{m} \gg \exp(c\sqrt{N})$. To prove (*) we need more advanced technique.

For an N -dimensional random vector X by $X_1^* \geq X_2^* \geq \dots \geq X_N^*$ we denote the nonincreasing rearrangement of $|X(1)|, \dots, |X(N)|$ (in particular $X_1^* = \max\{|X(1)|, \dots, |X(N)|\}$ and $X_N^* = \min\{|X(1)|, \dots, |X(N)|\}$). Random variables X_k^* , $1 \leq k \leq N$, are called order statistics of X .

Problem: Find an upper bound for $\mathbb{P}(X_k^* \geq t)$.

If coordinates of X_i are independent symmetric exponential r.v. with variance 1 then $\text{Med}(X_k^*) \sim \log(eN/k)$ for $k \leq N/2$.

Theorem

Let X be N -dim. log-concave isotropic vector. Then for all $t \geq C \log\left(\frac{eN}{k}\right)$,

$$\mathbb{P}\left(X_k^* \geq t\right) \leq \exp\left(-\frac{1}{C}\sqrt{kt}\right).$$

Actually we need a stronger theorem that uses the following important “weak parameter”

For a vector X in \mathbb{R}^N we define

$$\sigma_X(p) := \sup_{t \in S^{N-1}} (\mathbb{E}|\langle t, X \rangle|^p)^{1/p} \quad p \geq 2.$$

We also let σ_X^{-1} to be the left-inverse function.

$$\sigma_X^{-1}(s) := \sup \{t: \sigma_X(t) \leq s\}$$

Examples

- For isotropic log-concave vectors X , $\sigma_X(p) \leq p/\sqrt{2}$.
- For subgaussian vectors X , $\sigma_X(p) \leq C\sqrt{p}$.
- We say that an isotropic vector X is ψ_α if $\sigma_X(p) \leq Cp^{1/\alpha}$ (uniform distributions on suitable normalized B_p^N balls are ψ_α with $\alpha = \min(p, 2)$)

Theorem

For any N -dim. log-concave isotropic vector X , $\ell \geq 1$ and $t \geq C \log\left(\frac{eN}{\ell}\right)$,

$$\mathbb{P}\left(X_\ell^* \geq t\right) \leq \exp\left(-\sigma_X^{-1}\left(\frac{1}{C}t\sqrt{\ell}\right)\right)$$

This theorem implies uniform estimates for norms of projections.

Theorem

Let X be an isotropic log-concave vector in \mathbb{R}^N , and $m \leq N$. Then for any $t \geq 1$,

$$\mathbb{P} \left(\sup_{|I|=m} |P_I X| \geq Ct\sqrt{m} \log \left(\frac{eN}{m} \right) \right)$$

is less than or equal to

$$\exp \left(- \sigma_X^{-1} \left(\frac{t\sqrt{m} \log \left(\frac{eN}{m} \right)}{\sqrt{\log(em)}} \right) \right) \leq \exp \left(- t \frac{\sqrt{m}}{\sqrt{\log(em)}} \log \left(\frac{eN}{m} \right) \right).$$

The bound is sharp, except for $\sqrt{\log}$ in the probability estimate.
We conjecture this factor is not needed.

Idea of the proof. We want to estimate

$$\mathbb{P}\left(\sup_{|I|=m} |P_I X| \geq Ct\sqrt{m} \log\left(\frac{eN}{m}\right)\right).$$

We have

$$\sup_{|I|=m} |P_I X| = \left(\sum_{k=1}^m |X_k^*|^2\right)^{1/2} \leq 2\left(\sum_{i=0}^{s-1} 2^i |X_{2^i}^*|^2\right)^{1/2},$$

where $s = \lceil \log_2 m \rceil$.

We then use the estimates for order statistics.

Our approach to estimates for order statistics is based on the suitable estimate of moments of the process $N_X(t)$, where for $t \geq 0$,

$$N_X(t) := \sum_{i=1}^N \mathbb{1}_{\{X_i \geq t\}}$$

That is, $N_X(t)$ is equal to the number of coordinates of X larger than or equal to t .

Theorem

For any isotropic log-concave vector X and $p \geq 2$ we have

$$\mathbb{E}(t^2 N_X(t))^p \leq (Cp)^{2p} \quad \text{for } t \geq C \log \left(\frac{Nt^2}{p^2} \right).$$

$$\mathbb{E}(t^2 N_X(t))^p \leq (C\sigma_X(p))^{2p} \quad \text{for } t \geq C \log \left(\frac{Nt^2}{\sigma_X^2(p)} \right).$$

Estimate for N_X

Theorem

For any isotropic log-concave vector X and $p \geq 2$ we have

$$\mathbb{E}(t^2 N_X(t))^p \leq (Cp)^{2p} \quad \text{for } t \geq C \log \left(\frac{Nt^2}{p^2} \right).$$

$$\mathbb{E}(t^2 N_X(t))^p \leq (C\sigma_X(p))^{2p} \quad \text{for } t \geq C \log \left(\frac{Nt^2}{\sigma_X^2(p)} \right).$$

To get estimate for order statistics we observe that $X_k^* \geq t$ implies that $N_X(t) \geq k/2$ or $N_{-X}(t) \geq k/2$ and vector $-X$ is also isotropic and log-concave, and $\sigma_X = \sigma_{-X}$. Estimates for N_X and Chebyshev's inequality give

$$\mathbb{P}(X_k^* \geq t) \leq \left(\frac{2}{k} \right)^p (\mathbb{E}N_X(t)^p + \mathbb{E}N_{-X}(t)^p) \leq \left(\frac{C\sigma_X(p)}{t\sqrt{k}} \right)^{2p}$$

provided that $t \geq C \log(Nt^2/\sigma_X^2(p))$. Set $p = \sigma_X^{-1} \left(\frac{1}{eC} t\sqrt{k} \right)$ and notice that the restriction on t follows by the assumption that $t \geq C \log(eN/k)$.

Convolutions of measures

Let X_1, \dots, X_n be independent isotropic log-concave vectors in \mathbb{R}^N , and let $x = (x_i) \in \mathbb{R}^n$. Consider the vector

$$Y = \sum_{i=1}^n x_i X_i \in \mathbb{R}^N.$$

Probability for bounds of the process

$$\sup_{|I|=m} |P_I Y|$$

depends on the vector x . Specifically, it is convenient to assume the normalization $|x| \leq 1$ and $\|x\|_\infty \leq b \leq 1$.

Uniform bound for projections of convolutions

Theorem

Let $Y = \sum_{i=1}^n x_i X_i$, where X_1, \dots, X_n are independent isotropic N -dimensional log-concave vectors. Assume that $|x_i| \leq 1$ and $\|x\|_\infty \leq b \leq 1$.

i) If $b \geq \frac{1}{\sqrt{m}}$, then for any $t \geq 1$,

$$\mathbb{P} \left(\sup_{|I|=m} |P_I Y| \geq Ct\sqrt{m} \log \left(\frac{eN}{m} \right) \right) \leq \exp \left(- \frac{t\sqrt{m} \log \left(\frac{eN}{m} \right)}{b\sqrt{\log(e^2 b^2 m)}} \right).$$

ii) If $b \leq \frac{1}{\sqrt{m}}$ then for any $t \geq 1$,

$$\begin{aligned} \mathbb{P} \left(\sup_{|I|=m} |P_I Y| \geq Ct\sqrt{m} \log \left(\frac{eN}{m} \right) \right) \\ \leq \exp \left(- \min \left\{ t^2 m \log^2 \left(\frac{eN}{m} \right), \frac{t}{b} \sqrt{m} \log \left(\frac{eN}{m} \right) \right\} \right). \end{aligned}$$

Uniform bound for norms of submatrices

Let A be an $n \times N$ random matrix with independent log-concave isotropic rows $X_1, \dots, X_n \in \mathbb{R}^N$. For $k \leq n, m \leq N$ we let

$A_{k,m}$ = the maximal operator norm of a $k \times m$ submatrix of A .

For simplicity assume $n \leq N$.

Theorem

For any integers $n \leq N, k \leq n, m \leq N$ and any $t \geq 1$, we have

$$\mathbb{P}\left(A_{k,m} \geq Ct\lambda_{mk}\right) \leq \exp\left(-\frac{t\lambda_{mk}}{\sqrt{\log(3m)}}\right),$$

where

$$\lambda_{mk} = \sqrt{\log \log(3m)} \sqrt{m} \log\left(\frac{eN}{m}\right) + \sqrt{k} \log\left(\frac{en}{k}\right).$$

The threshold value of λ_{km} is optimal, up to the factor $\sqrt{\log \log(3m)}$. Under the additional assumption of unconditionality of the rows we can remove this factor and get the sharp estimate. For example, for expectations:

Theorem

Let A be an $n \times N$ matrix whose rows are independent isotropic log-concave unconditional random vectors. Then

$$\mathbb{E}A_{km} \leq C \left(\sqrt{m} \log \frac{3N}{m} + \sqrt{k} \log \frac{3n}{k} \right).$$

Reconstruction

Let n, N be integers; $X \in \mathbb{R}^N$ a random vector, X_1, \dots, X_n independent, distributed as X . A is an $n \times N$ matrix which has X_i 's as rows

$$A = \begin{pmatrix} X_1 \\ X_2 \\ \dots \\ X_n \end{pmatrix} = \begin{pmatrix} \dots\dots\dots \\ \dots\dots\dots \\ \dots\dots\dots \\ \dots\dots\dots \end{pmatrix}$$

We can treat $A : \mathbb{R}^N \rightarrow \mathbb{R}^n$ given by $x \mapsto (\langle X_i, x \rangle) \in \mathbb{R}^n$.

Let $m \leq N$. A vector $x \in \mathbb{R}^N$ is m -sparse if $|\text{supp } x| \leq m$.

Problem from Compressed Sensing theory: Reconstruct any m -sparse vector $x \in \mathbb{R}^N$ from the data $Ax \in \mathbb{R}^n$, with a fast algorithm.

Given Ax , find x , knowing that it is sparse. Note that of course A is not-invertible.

RIP and Geometry of Polytopes

Define $\delta_m = \delta_m(A)$ as the infimum of $\delta > 0$ such that

$$||Ax|^2 - \mathbb{E}|Ax|^2| \leq \delta n |x|^2$$

holds for all m -sparse vectors $x \in \mathbb{R}^N$.

δ_m is the Restricted Isometry Property (RIP) parameter of order m .

Candes and Tao (2006): if δ_{2m} is sufficiently small then

- (*) whenever $y = Ax$ has a m -sparse solution x , then x is the unique solution of the ℓ_1 -minimization program: $\min ||t||_{\ell_1}$ with the min over all t such that $At = y$.

Geometry of Polytopes: By Donoho (2005), (*) is equivalent to the condition that the centrally symmetric polytope $A(B_1^N)$ is m -centrally neighborly (i.e., any set of less than m vertices containing no opposite pairs, is a vertex set of a face).

Lemma

Let X_1, \dots, X_n be independent isotropic random vectors in \mathbb{R}^N . Let $0 < \theta < 1$ and $B \geq 1$. Then with probability at least

$$1 - \binom{N}{m} \exp(-3\theta^2 n / 8B^2)$$

one has

$$\delta_m \leq \theta + \frac{1}{n} (A_{k,m}^2 + \mathbb{E}A_{k,m}^2),$$

where $k \leq n$ is the largest integer satisfying $k \leq (A_{k,m}/B)^2$;

RIP theorem for matrices with independent rows:

Theorem

Let $n \leq N$ and $0 < \theta < 1$. Let A be an $n \times N$ matrix, whose rows are independent isotropic log-concave random vectors X_i , $i \leq n$.

There exists an absolute constant $c > 0$, such that if $m \leq N$ satisfies

$$m \log \log(3m) \left(\log \frac{3N}{m} \right)^2 \leq c \left(\frac{\theta}{\log(3/\theta)} \right)^2 n$$

then

$$\delta_m \leq \theta$$

with overwhelming probability.

Optimal up to a log log factor.

For unconditional distributions can be removed