

Monotone cellular automata

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(Based on joint work with Paul Balister, József Balogh, Béla Bollobás, Hugo Duminil-Copin, Ivailo Hartarsky, Fabio Martinelli, Paul Smith, and Cristina Toninelli.)

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Note that this is a random variable, and is also a function of the initial state, and of p . An interesting particular case is when the initial state is chosen randomly (e.g., with density p of empty sites), and $p \rightarrow 0$.

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Let $A = A_0$ denote the set of initially infected sites, and define

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for each $t \geq 0$.

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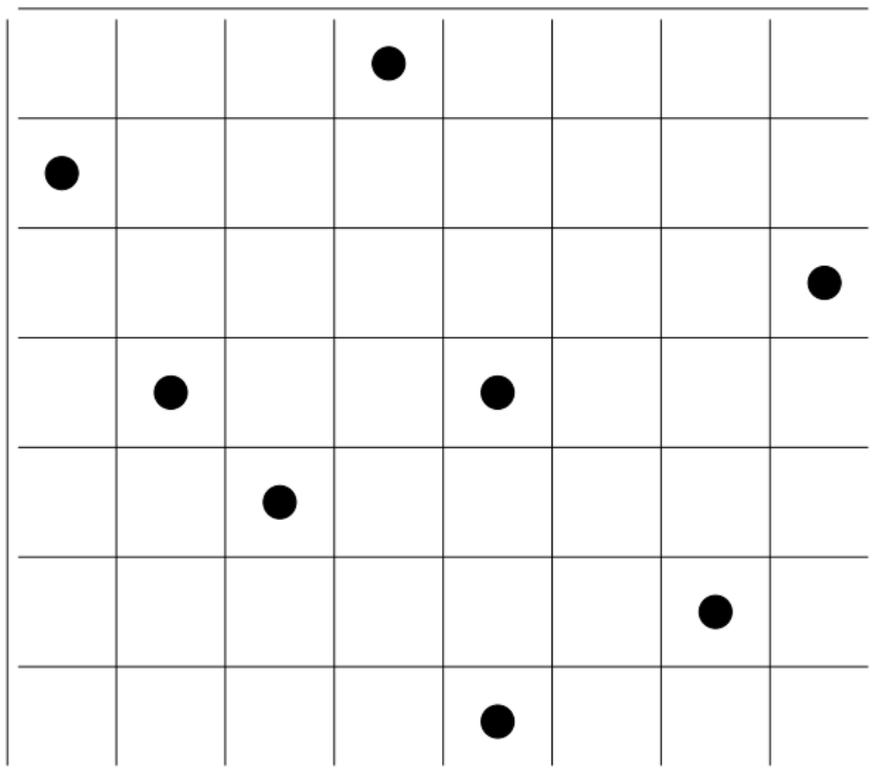
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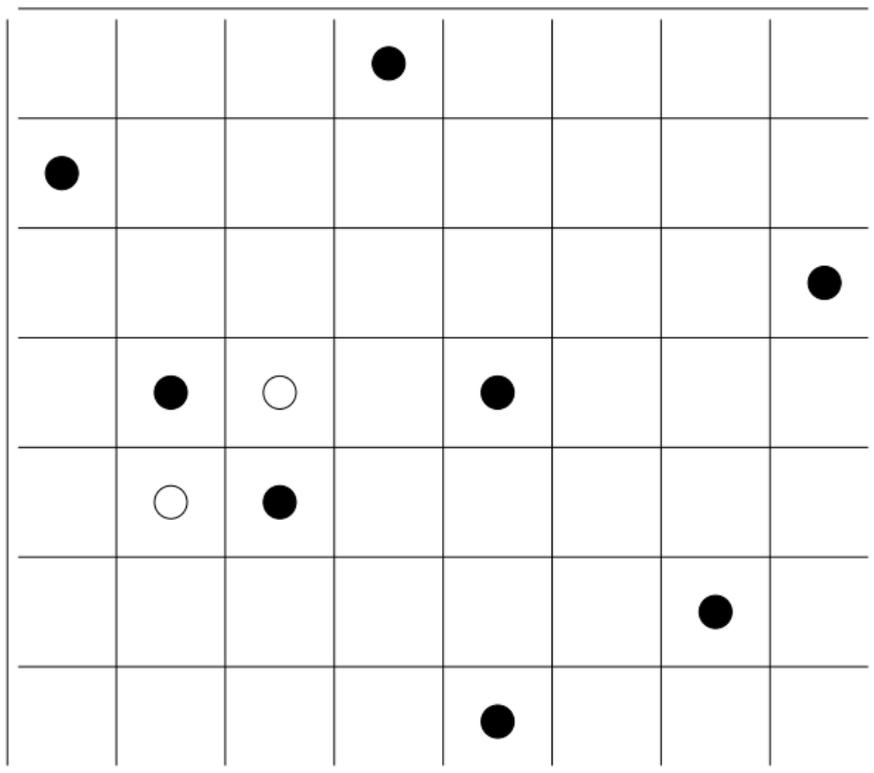
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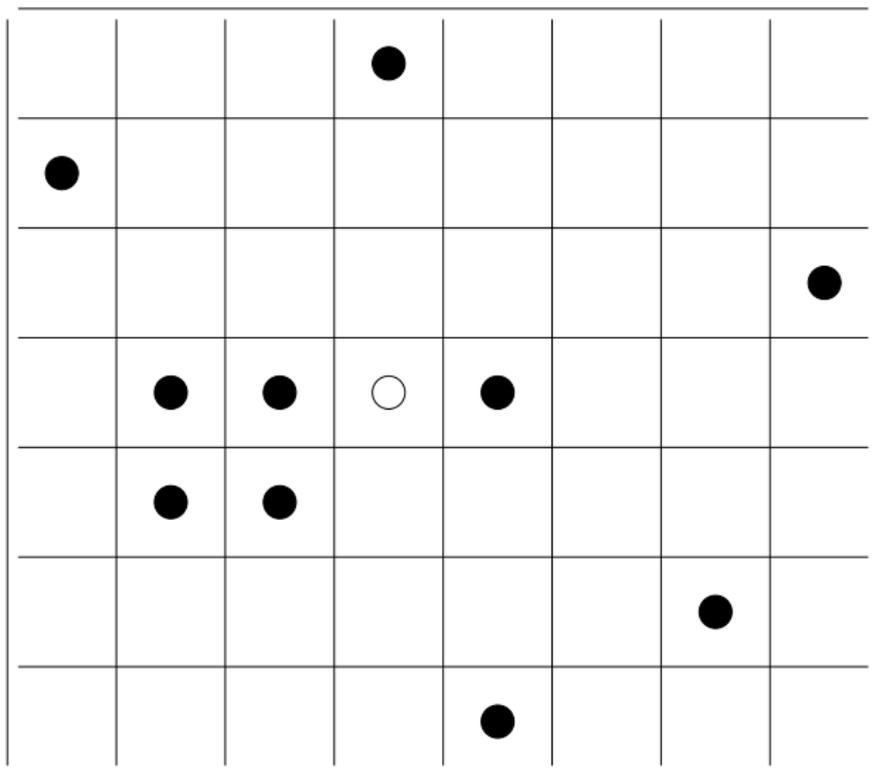
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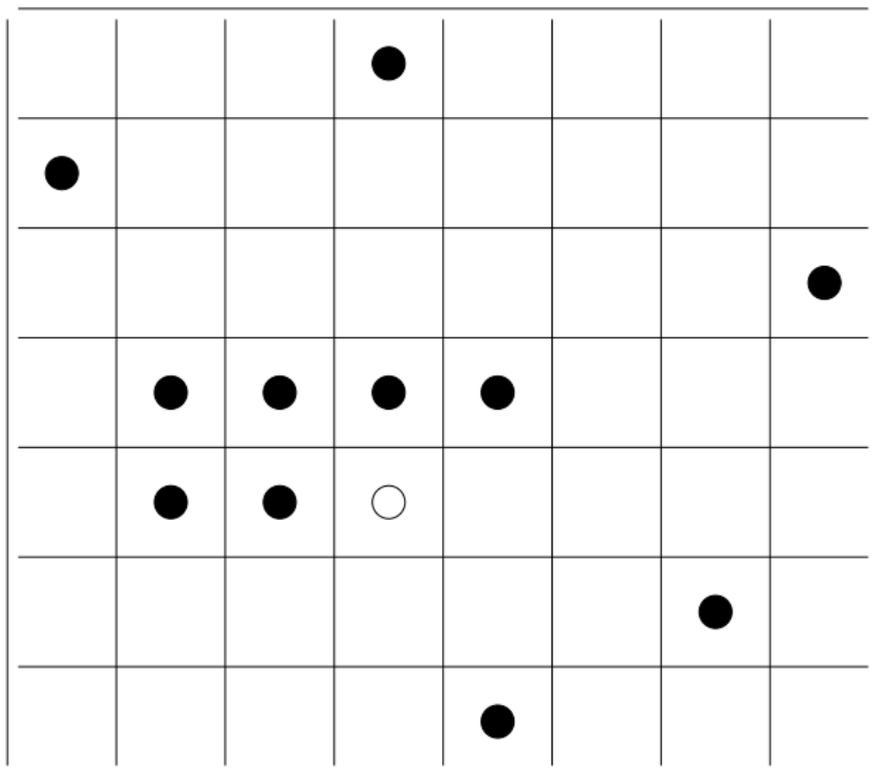
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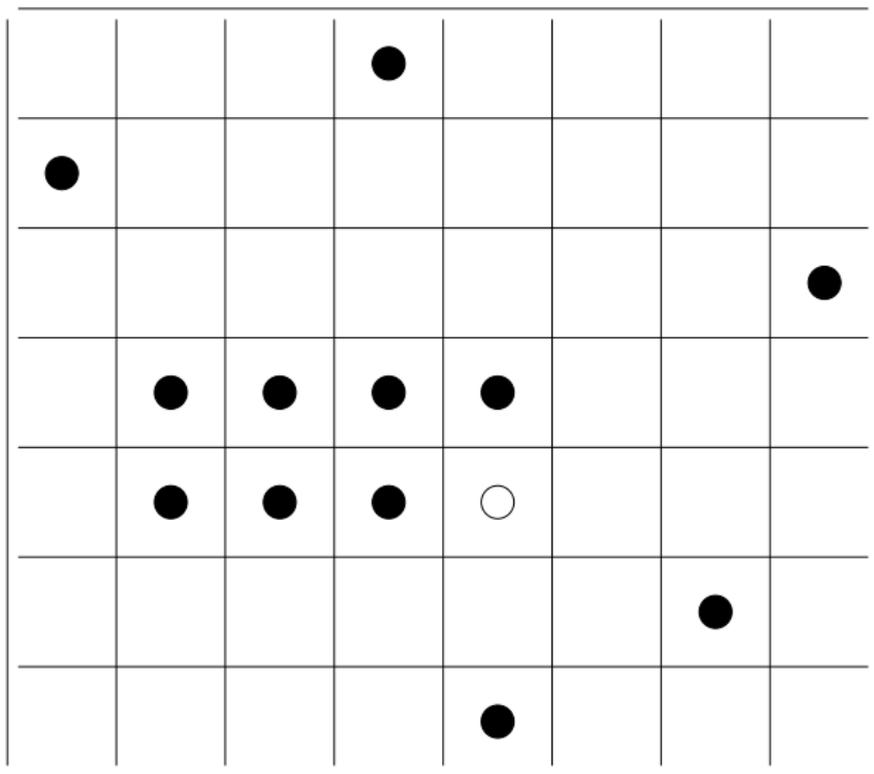
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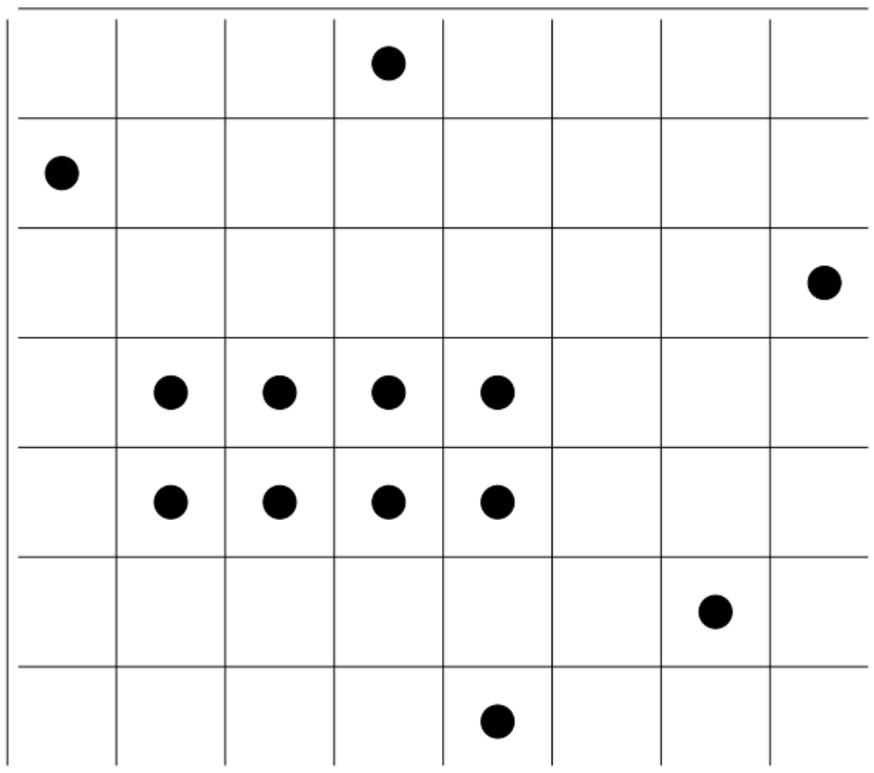
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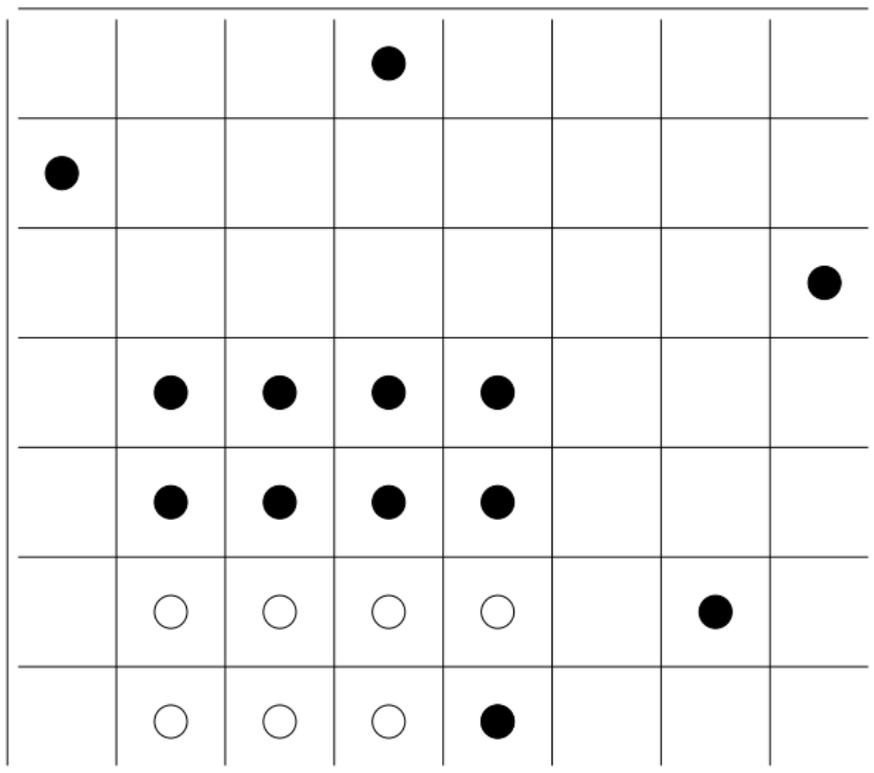
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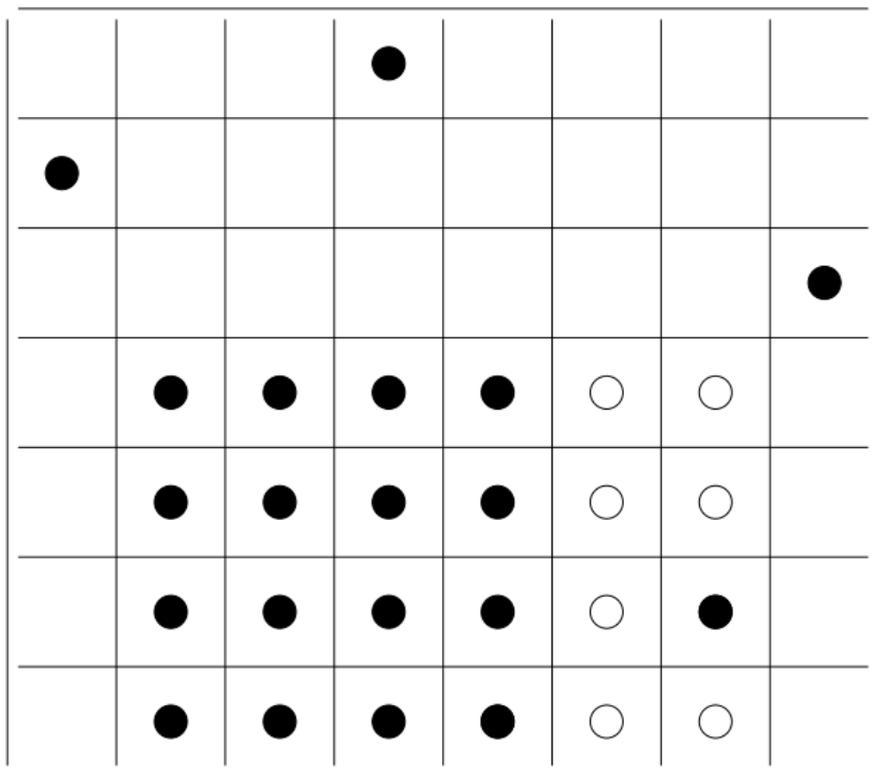
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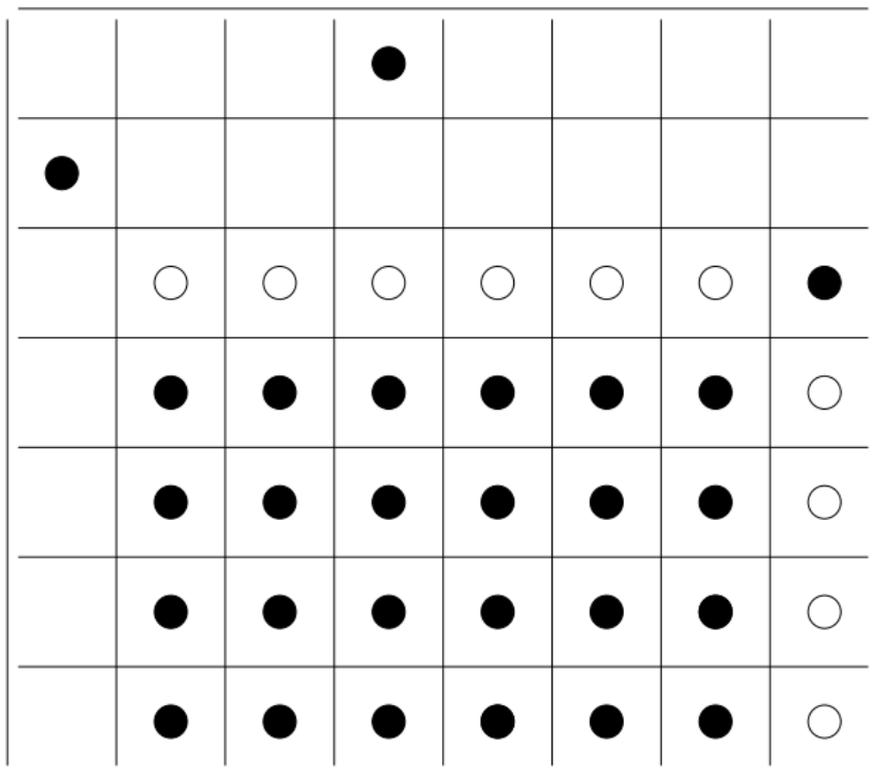
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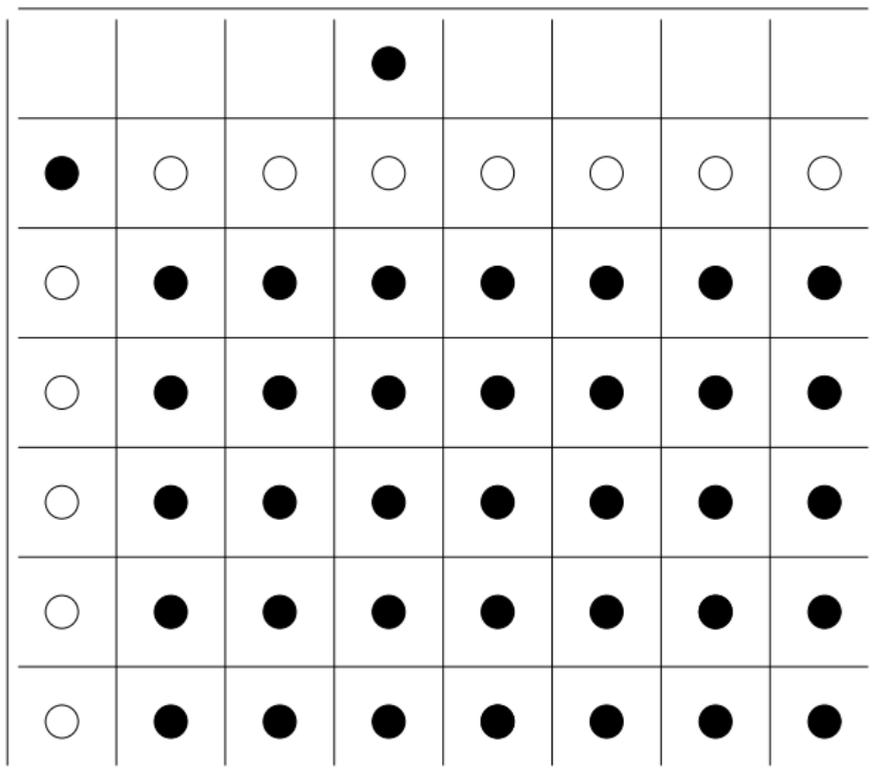
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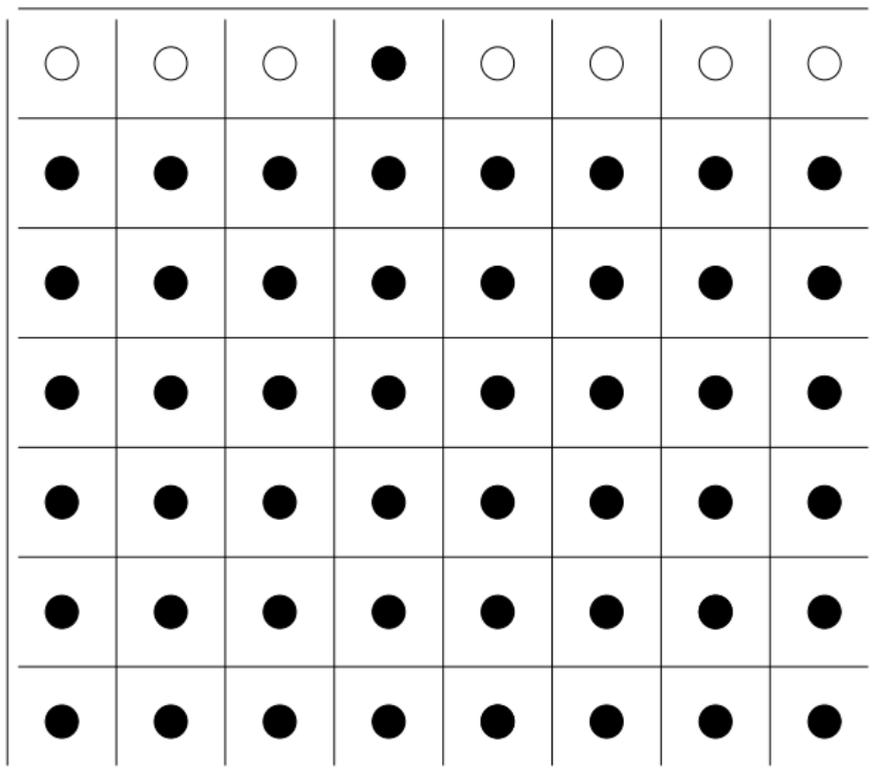
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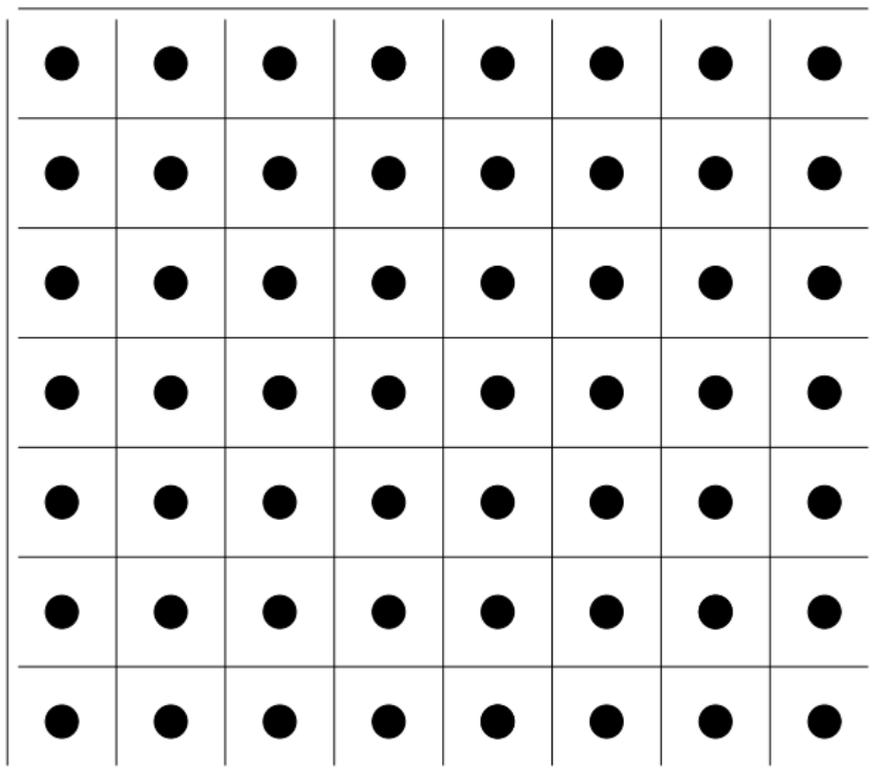
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Answer: $p_c(\mathbb{Z}^2, 2) = 0$ (!!)

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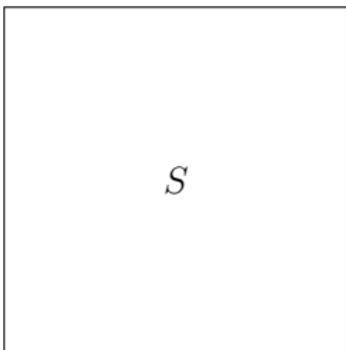
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Answer: $p_c(\mathbb{Z}^2, 2) = 0$ (van Enter, 1987)

van Enter's proof that $p_c(\mathbb{Z}^2, 2) = 0$ (sketch)

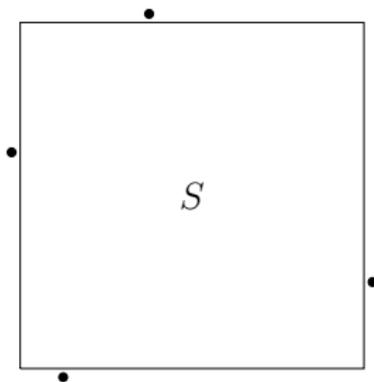
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With probability 1, there exists a very large completely infected square S (a *critical droplet*) somewhere in \mathbb{Z}^2 :



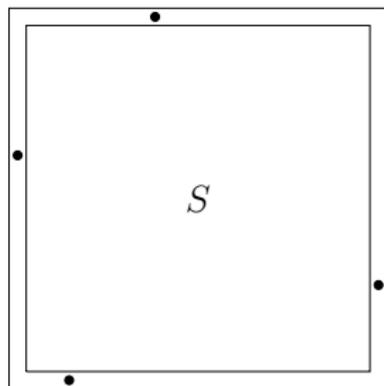
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Since S is very large, it is likely to have infected sites on its sides, and hence to be able to grow by one in each direction:



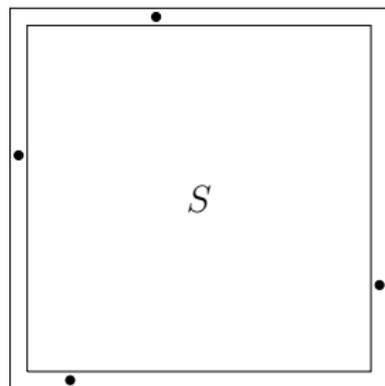
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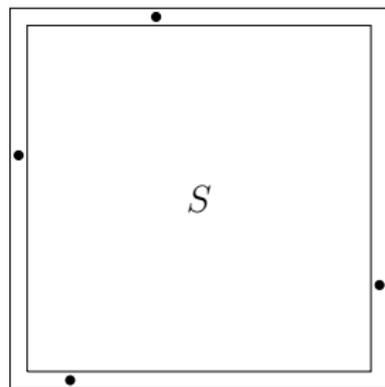
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We define the *critical probability* on an $n \times n$ torus to be

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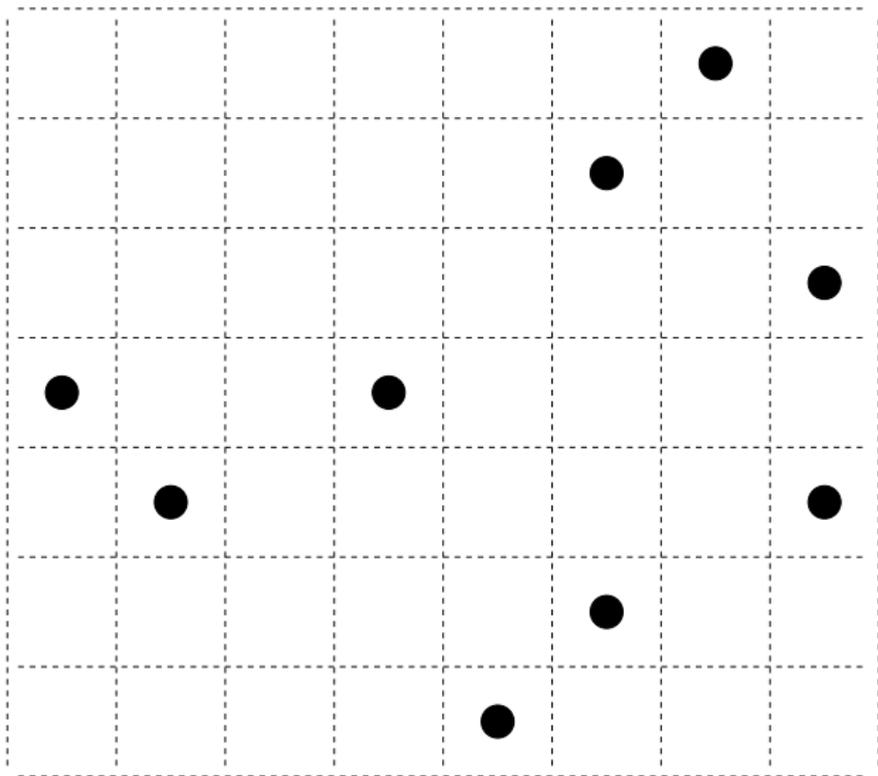
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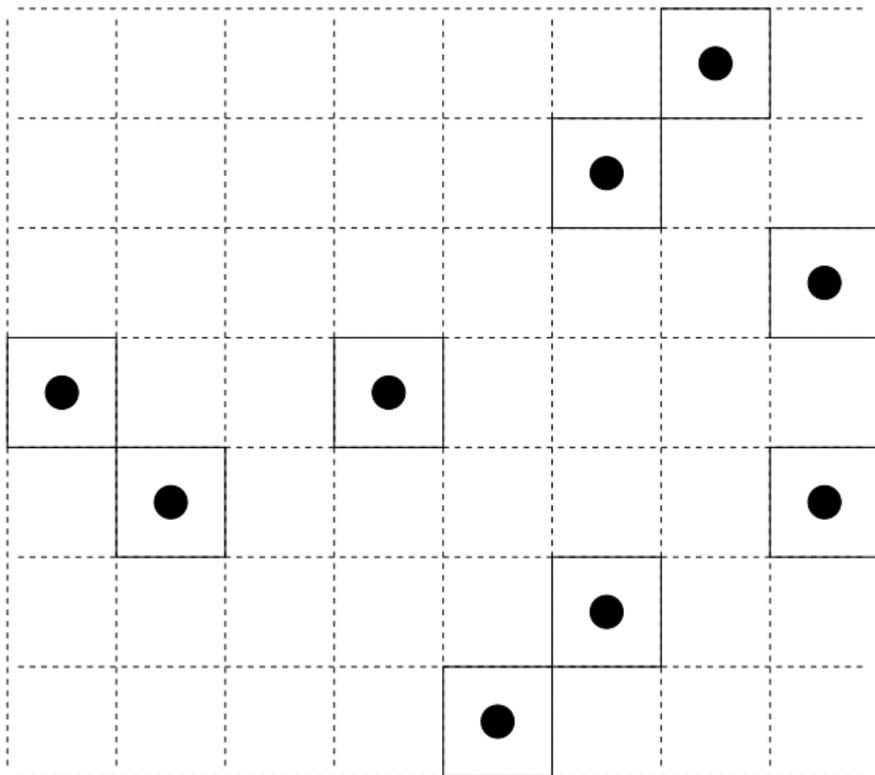
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- Every rectangle R that appears at some point in the rectangles process is *internally filled* by A , i.e., $[A \cap R] = R$.

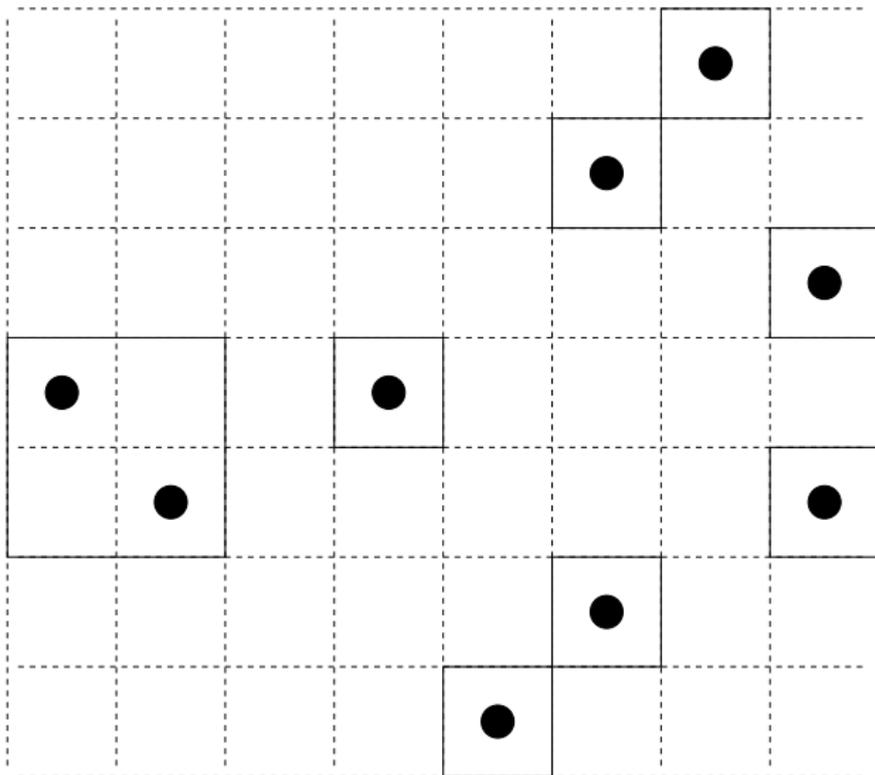
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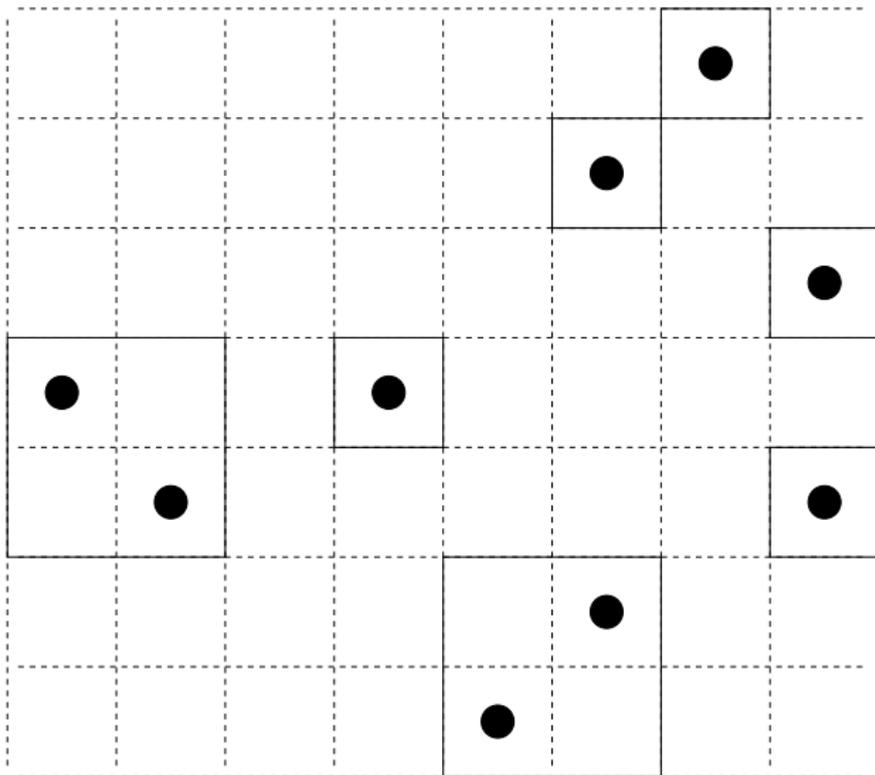
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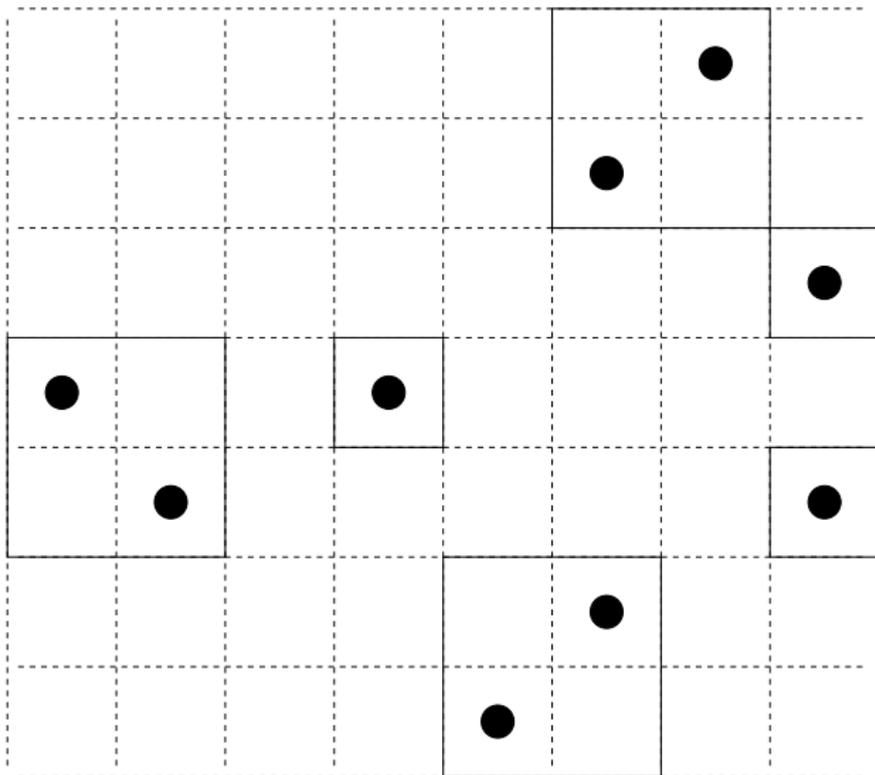
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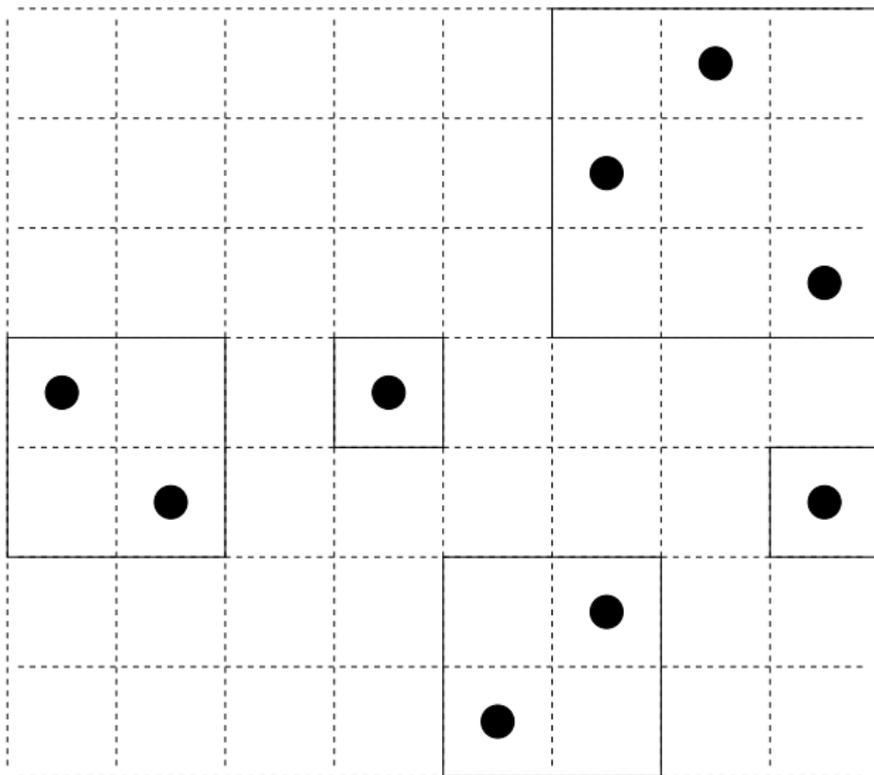
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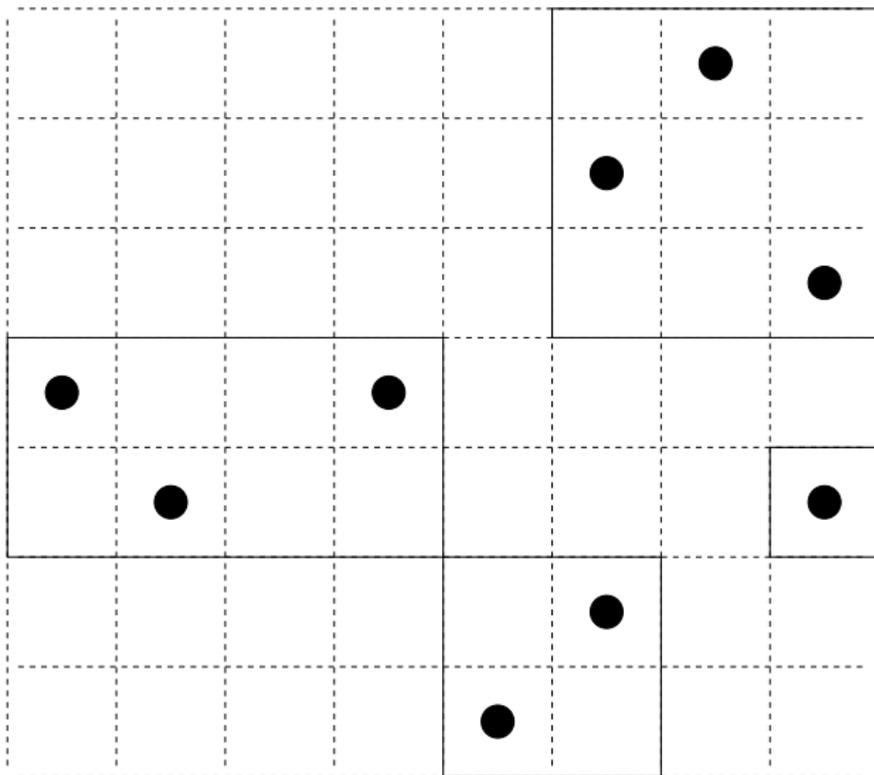
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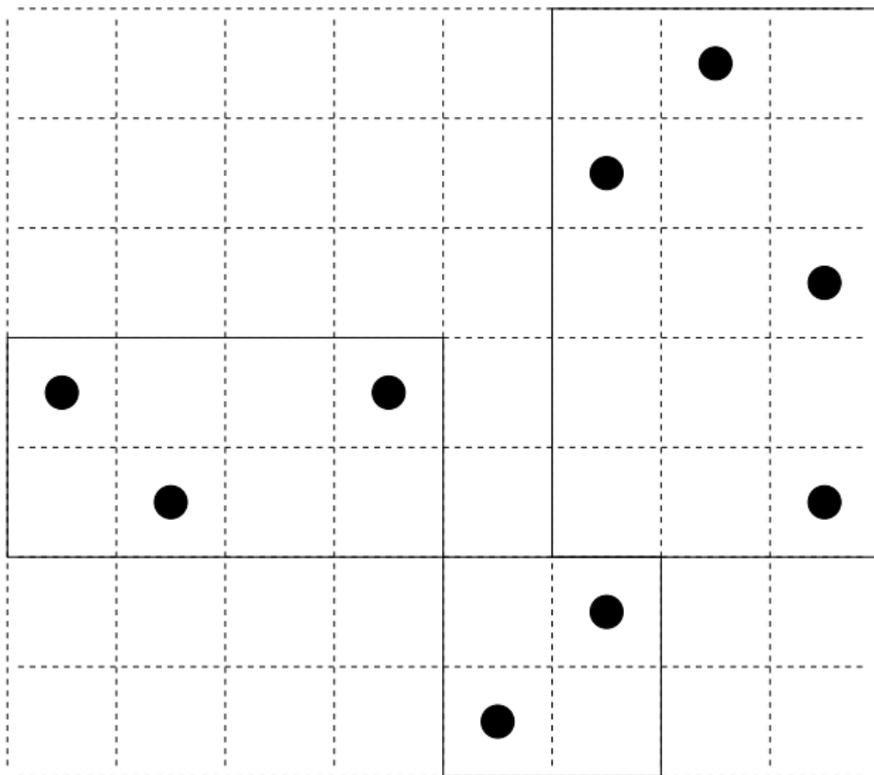
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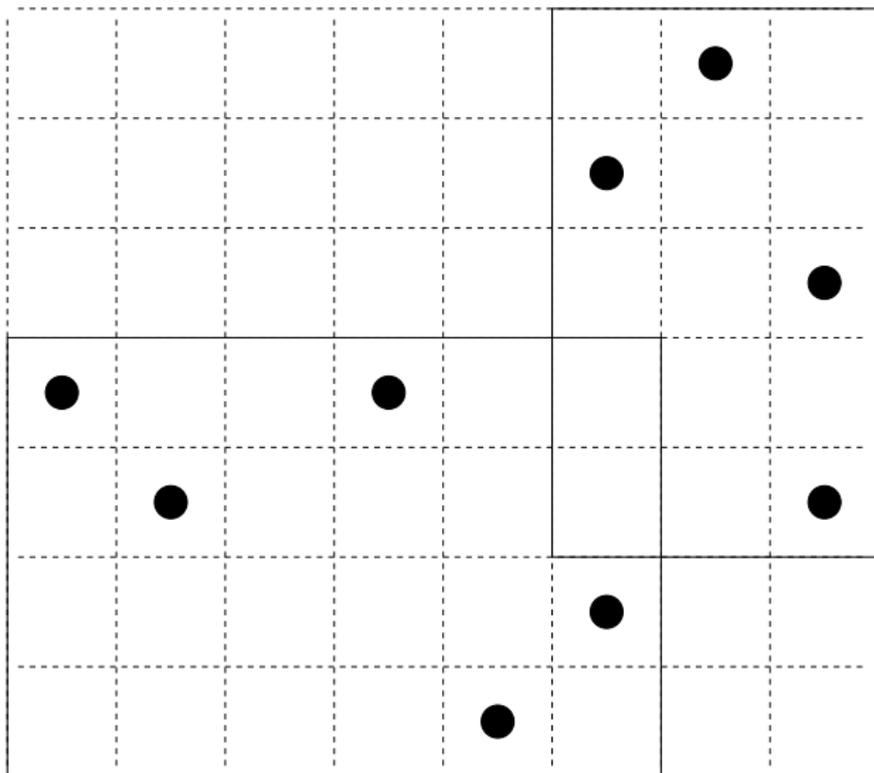
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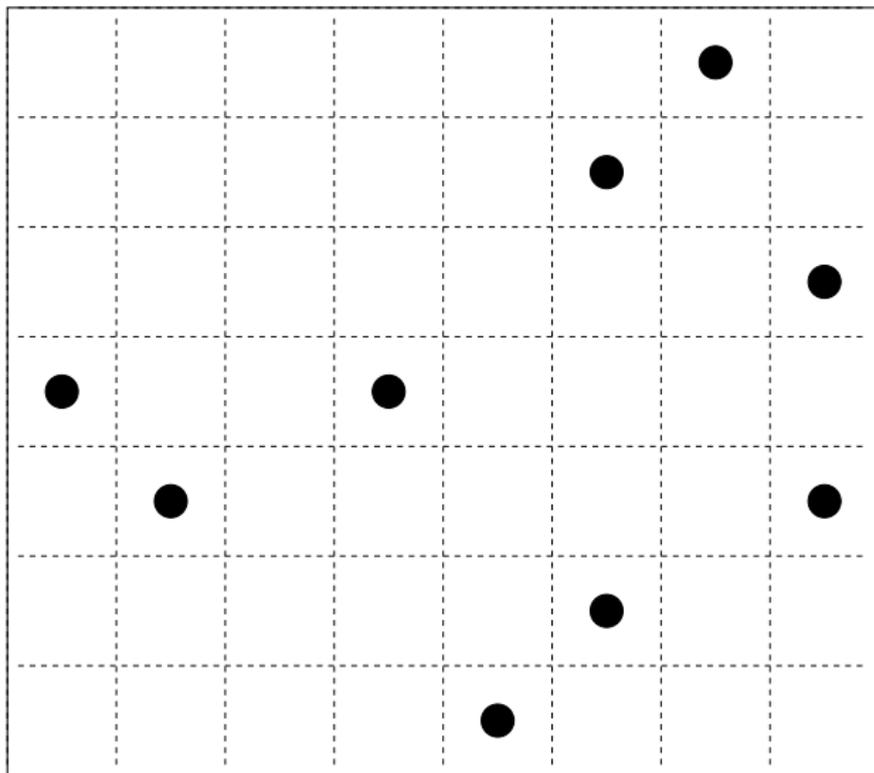
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Aizenman and Lebowitz's proof (sketch, continued)

Using the rectangles process, we can prove the following key lemma.

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The Aizenman–Lebowitz Lemma

If A percolates in \mathbb{Z}_n^2 , then there exists a rectangle $R \subset \mathbb{Z}_n^2$, with

$$\log n \leq \text{long}(R) \leq 2 \log n,$$

that is “internally filled”, i.e., $[A \cap R] = R$.

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Proof: Run the rectangles process until a rectangle with $\text{long}(R) \geq \log n$ appears for the first time. This rectangle is internally filled, by the definition of the process. Moreover, it was obtained from two rectangles with $\text{long}(R) < \log n$, so we have $\text{long}(R) \leq 2 \log n$, as required.

Aizenman and Lebowitz's proof (sketch, final calculation)

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$$p = \frac{\varepsilon}{\log n}$$

for some small constant $\varepsilon > 0$, then we obtain

$$\mathbb{P}([A \cap R] = R) \leq (4p \log n)^{\log n/2} \leq (4\varepsilon)^{\log n/2} \leq \frac{1}{n^3}.$$

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There are $n^3(\log n)^{O(1)}$ choices for R , so by Markov's inequality

$$\mathbb{P}(A \text{ percolates}) \rightarrow 0$$

as $n \rightarrow \infty$, as required.

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Recall that the states of sites at time zero are chosen independently at random, with density p of $+s$, and when a clock rings a site updates to agree with the majority of its four nearest neighbours; if it has two neighbours in each state, then it chooses a new state randomly.

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Conjecture (Folklore)

If $p > 1/2$ then the system fixates.

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If $p > 1 - 10^{-10}$ then the system fixates.

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Combining the proof of this theorem with some more advanced techniques from bootstrap percolation (see Balogh, Bollobás and M., 2009) one can prove the following result in high dimensions.

Theorem (M., 2011)

If $p > \frac{1}{2}$ and $d \geq d_0(p)$, then on \mathbb{Z}^d the system fixates.

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Theorem (Martinelli and Toninelli, 2017+)

There exist constants $C > c > 0$ such that

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with high probability as $p \rightarrow 0$.

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Sharp thresholds and higher dimensions

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Theorem (Aizenman and Lebowitz, 1988)

$$p_c(\mathbb{Z}_n^2, 2) = \frac{\Theta(1)}{\log n}.$$

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Theorem (Holroyd, 2003)

$$p_c(\mathbb{Z}_n^2, 2) = \left(\frac{\pi^2}{18} + o(1) \right) \frac{1}{\log n}.$$

Sharp thresholds and higher dimensions

For the 2-neighbour bootstrap model, much more precise bounds are known. Gravner and Holroyd later refined the upper bound argument, and together with them we proved an almost matching lower bound:

Theorem (Gravner–Holroyd, 2008; Gravner–Holroyd–M., 2012)

There exist constants $C > c > 0$ such that

$$\frac{\pi^2}{18 \log n} - \frac{C(\log \log n)^3}{(\log n)^{3/2}} \leq p_c(\mathbb{Z}_n^2, 2) \leq \frac{\pi^2}{18 \log n} - \frac{c}{(\log n)^{3/2}}$$

for every sufficiently large $n \in \mathbb{N}$.

Sharp thresholds and higher dimensions

For the 2-neighbour bootstrap model, much more precise bounds are known. Finally, with Hartarsky, we have managed to determine the order of the second term:

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The proof of Aizenman and Lebowitz also works in higher dimensions, but only for the 2-neighbour model:

Theorem (Aizenman and Lebowitz, 1988)

$$p_c(\mathbb{Z}_n^d, 2) = \left(\frac{\Theta(1)}{\log n} \right)^{d-1}.$$

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For the 3-neighbour model in three dimensions, the threshold was determined up to a constant factor by Cerf and Cirillo:

Theorem (Cerf and Cirillo, 1999)

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Theorem (Balogh, Bollobás and M., 2009)

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For the r -neighbour model in d dimensions, the threshold was determined up to a constant factor by Cerf and Manzo:

Theorem (Cerf and Manzo, 2002)

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Theorem (Balogh, Bollobás, Duminil-Copin and M., 2012)

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The Bollobás–Smith–Uzzell model

We now turn our attention to some dramatic recent developments in the study of bootstrap percolation, which were initiated a few years ago in a remarkable paper of Béla Bollobás, Paul Smith, and Andrew Uzzell. They studied the following large family of models:

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Definition (The \mathcal{U} -bootstrap process)

Let $\mathcal{U} = \{X_1, \dots, X_m\}$ be an arbitrary finite collection of finite subsets of \mathbb{Z}^2 , and let $A \subset \mathbb{Z}_n^2$. Set $A_0 = A$, and define, for each $t \geq 0$,

$$A_{t+1} = A_t \cup \left\{ x \in \mathbb{Z}_n^2 : x + X \subset A_t \text{ for some } X \in \mathcal{U} \right\}.$$

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One of the key insights of Bollobás, Smith and Uzzell was that the typical global behaviour of the \mathcal{U} -bootstrap process (with random initial set) should be determined by the action of the process on discrete half-spaces.

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Define $\mathcal{S} = \mathcal{S}(\mathcal{U}) \subseteq S^1$, the collection of *stable directions*, to be the set

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Let \mathcal{C} denote the collection of open semicircles in S^1 . The following key definition is due to Bollobás, Smith and Uzzell:

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Note that this is a partition of the two-dimensional update families.

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The first two parts of the following theorem were proved by Bollobás, Smith and Uzzell; the proof for subcritical families was obtained slightly later by Balister, Bollobás, Przykucki and Smith.

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The threshold for critical models

With Bollobás, Duminil-Copin and Smith, we proved the following much more precise bounds for critical update families:

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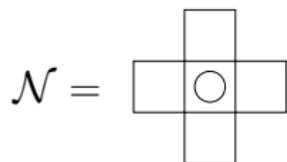
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\mathcal{U} is balanced if and only if there exists a *closed* semicircle such that $\alpha(u) \leq \alpha$ for every $u \in C$.

Critical models: some examples

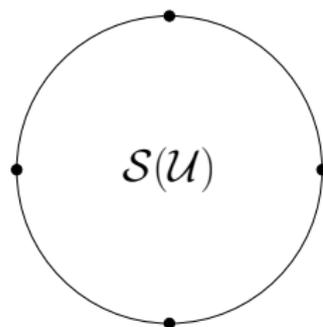
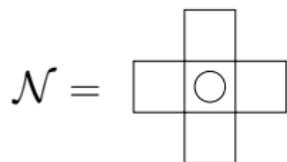
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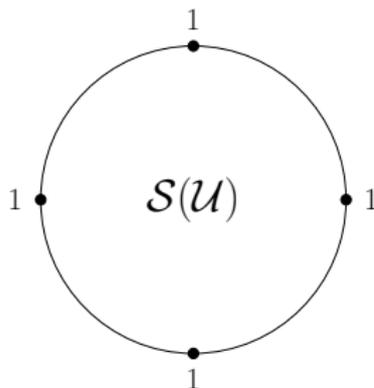
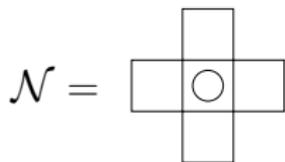
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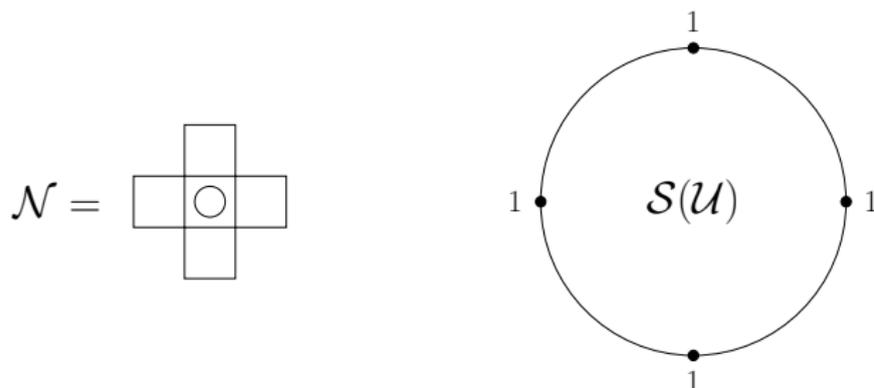
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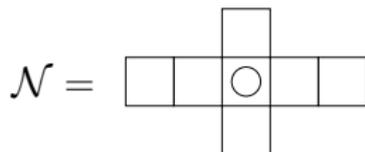
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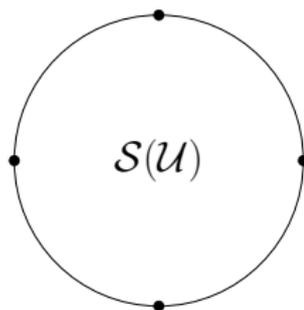
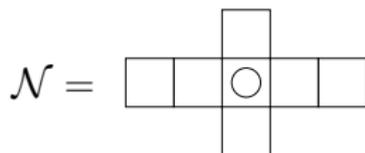
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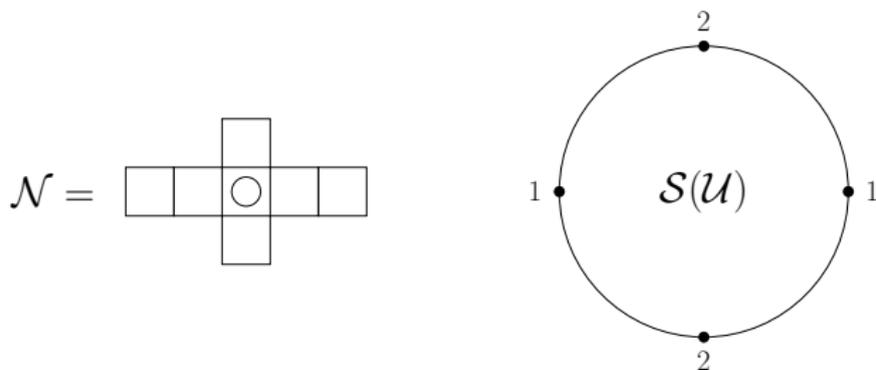
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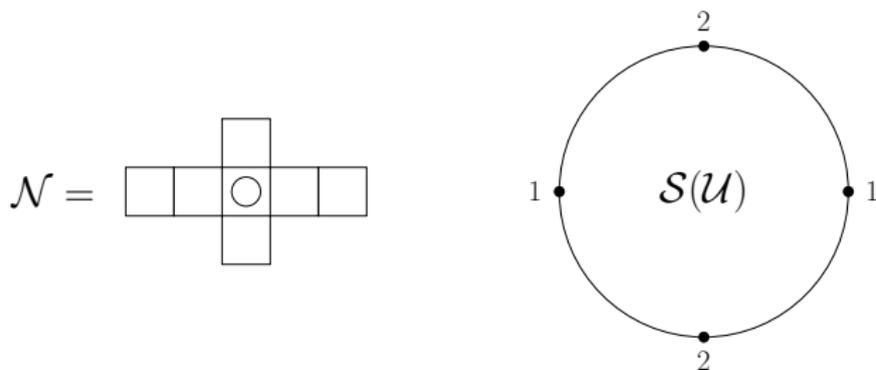
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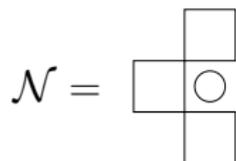
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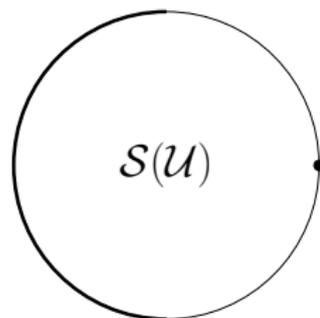
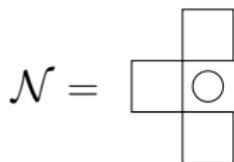
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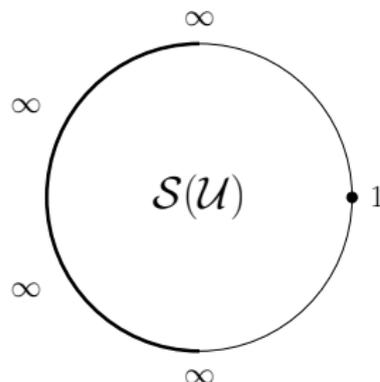
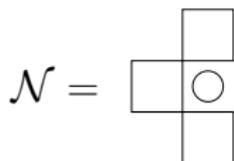
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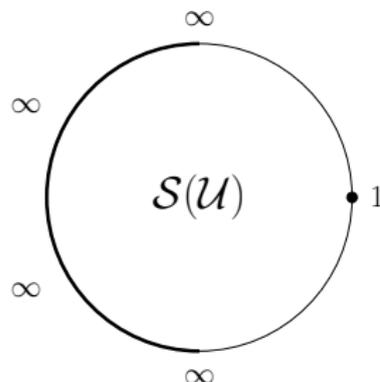
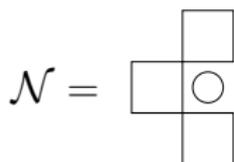
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For every critical unrooted update family \mathcal{U} ,

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Conjecture (Martinelli, M. and Toninelli, 2017+)

For every critical rooted update family \mathcal{U} , there exists $\beta > \alpha$ such that

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Thank you!

Universality for higher dimensions

Theorem (Balister–Bollobás–M.–Smith, 2017+)

Let \mathcal{U} be a d -dimensional update family.

(a) If \mathcal{U} is supercritical then $p_c(\mathbb{Z}_n^d, \mathcal{U}) = n^{-\Theta(1)}$,

(b) If \mathcal{U} is critical then there exists $r = r(\mathcal{U}) \in \{2, \dots, d\}$ such that

$$p_c(\mathbb{Z}_n^d, \mathcal{U}) = \left(\frac{1}{\log_{(r-1)} n} \right)^{\Theta(1)},$$

(c) If \mathcal{U} is subcritical then $p_c(\mathbb{Z}^d, \mathcal{U}) > 0$.

When $r < d$, the constant in the power is in general uncomputable (!!)